ALICE

CERN, SCCS Technische Universität München Fast and Efficient Entropy Compression of ALICE Data using ANS Coding

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OBJECTIV

• Provide state of the art lossles entropy compression in the ALICE O² framework that minimizes permanent storage requirements of detector data in LHC run 3.

- Preprocessed detector data are provided as *Compressed Time Frames* as flat, numeral arrays to be written out as a binary stream to permanent storage at a target data rate of ~90GB/s [1]
- Coder must efficienlty deal with source alphabets of up to 25 bits per symbol.
- If possible, exploit CPU and GPU resources of heterogenious ALICE O² farm.
- Maintain compatibility between data created by CPU and GPU.

ALICE TIME PROJECTION CHAMBER (TPC) DATA

• Alice TPC responsible for >90% of data in ALICE timeframe for LHC run 3 [1].

ALICE Timeframe

SAMPLE SOURCE DATA

- Flat structure of arrays, interpretable as four tables with multiple columns.
- Source symbol value range of 8-25 Bits unsigned integers per column.
- *Variable range coders* use symbol distribution of source data for compression.
- Examples for Variable range coders are Huffmann Coding or Arithmetic Coding.



- ANS by Jarek Duda [2] adds new family of *variable range coders*.
- rANS as fast entropy coder for large alphabets based on arithmetic operations.
- Encodes symbols into single, infinite precision integer state in LiFo order.
- Decoder unwinds LiFo stack, reversing coder steps.

Message: BAACAB, length 6 =: m



- 130 simulated Pb-Pb collisions provide $O(10^7)$ symbols for good statics.
- Non standard, patchy distribution of source symbols.



- Cross checks with detector data from LHC run 2 show high similarity.
- Assume static distribution of source symbols within each column.
- Calculation of entropy (H) as lower bound for compression.
- Concatenation of symbols in correlated columns to reduce overall entropy.

| • Calculation of entropy (n) as lower bound for compression. | | | | | |
|---|----------|------|----------|--------------|---------------------|
| Concatenation of symbols in correlated columns to reduce overall entropy. | | | | | |
| Table | Rows | Cols | Bits/row | H [Bits/row] | H concat [Bits/row] |
| AttachedClusters | 21072849 | 4 | 41 | 17.15 | 15.75 |
| AttachedClustersReduced | 20590430 | 4 | 55 | 17.60 | 17.59 |
| Tracks | 482 419 | 5 | 73 | 53.90 | 53.90 |
| UnattachedClusters | 50745911 | 5 | 81 | 39.77 | 38.37 |
| | | | | | |

ENTROPY CODING TESTSBED



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ASYMETRI

 \bigcap

Ζ

• State interpretable as a cyclic, *asymetrized base-m numeral system*.



- State grows with probability of encoded symbol $x_{n+1} \approx x_n / Pr[s]$.
- overcomes 1Bit/symbol limitation of Huffman, on pair with Arithmetic Coding.
- Constrain x to a finite Interval I by streaming in/out bits. Streaming Coder



PERFORMANCE CONSIDERATIONS

- Rescaling $m \coloneqq 2^p$ allows mult, div be replaced by shifts and mod by bitops.
- Replace divisions by multiplication with precalculated inverse.
- rANS en-/decoding faster then Arithmetic Coding, slower then Huffman Coding. RANS AND PARALLELISM
- \bigcirc tion system. \leq Coded data is decoded af-Preparation Check Encode Decode terwards and compared to J Correctnes Checks source data for correctnes. フ • C++ rANS code based on reengineered *ryg_rans* [3] coder/decoder. S • Bandwidth: Single thread, (Bits/row×Rows) /time, avg. over 5 runs. S Ζ ATTAINABLE LEVEL OF ENTROPY COMPRESSION **G** Bits/row H [Bits/row] rANS [Bits/row] rANS/H BW [MiB/s] Table AttachedClusters 41 15.75 15.75 1.00 649.72 \triangleright **AttachedClustersReduced** 55 776.06 17.59 17.64 1.00 Tracks 73 53.90 53.90 1.00 365.56 \bigcirc UnattachedClusters 81 38.37 38.37 1.00 589.80 • rANS coder is close to entropy for all tables at high encoding speeds. **COMPRESSION AND LUT PRECISION** Rescaling LUT to 2^p for perfromance. \bigcirc - 16Bit Symbols LUT must map source distribution 2 1.30 25Bit Symbols DATA well enough for good compression. <u></u> 2 1.25 Trivialy: |LUT|>|SourceAlphabet|, i.e 1.20 larger LUTs for larger alphabets. . 1.15 ₹ 1.10 ‡ Must be preserved for decoding. S 1.05 Additional Inverse LUT (iLUT) for decoder |iLUT|≥|LUT|, constructed from 18 20 22 24 26 28 10 12 14 16 LUT size [log₂] LUT.
 - For large source alphabets only sensible for static distributions, where LUTs can be reused or for extremely large datasets.
 - Requires algorithmic improvements to decrease memory requirements.

- Coder and Decoder are exact inverse functions: D(C(s,x))=(s,x).
- Allows to write incompressible bits or parallel coders on same stream, see [4].
- Still requires operations by C(s,x) to remain exactly reversible by D(s).
- rANS is a powerful entropy coding algorithm for large alphabets.
- Capable of handling close to optimal compression of all tested ALICE datasets.
- Fast encoding performance, currently lacks in decoding with large alphabets.

OUTLOOK

ONCLUSION

References:

- Improve code robustness and decoding speeds by optimizing LUT.
- Finish Implementation of SIMD and GPU features.
- Integration into ALICE O².

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BANDWIDTH MEASUREMENTS





• Performance bound by lookups in

iLUT that does not fit into cache.

[3] Fabian Giesen, https://github.com/rygorous/ryg_rans, last checked 10/2019 [4] Fabian Giesen, Interleaved entropy coders, arXiv:1402.3392, 02/2014

+ 8Bit Symbols 16Bit Symbols

Only raw (de)coding speed is

measured. LUTs are staticaly

kept in memory for produc-

Rescale LUT